

Local energy minima of the 3d Edwards-Anderson model

Stefan Schnabel and Wolfhard Janke

CompPhys 16

Edwards-Anderson model

- Hamiltonian : $\mathcal{H}(\mathbf{S}_1, \dots, \mathbf{S}_N) = - \sum_{\langle ij \rangle} J_{ij} \mathbf{S}_i \mathbf{S}_j, \quad \mathbf{S} \in \{-1, 1\}$
- Gaussian disorder : $P(J) = \frac{1}{\sqrt{2\pi}} e^{-\frac{J^2}{2}}$
- Overlap of two configurations $Q(\mathbf{S}^\alpha, \mathbf{S}^\beta) = \frac{1}{N} \sum_i S_i^\alpha S_i^\beta$
- 3d – simple cubic lattice

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Greedy algorithm

- Reduces E by **locally** optimal steps.
- Find spin with the highest positive energy e_i , flip it and repeat until all $e_i < 0$.
- “Traditional” application : Randomize many configurations and minimize with greedy algorithm.
- Does not find the global minimum.
- Constitutes basins of attraction:

$$B(\mathbf{S}_{\min}) := \{\mathbf{S}_B : G(\mathbf{S}_B) = \mathbf{S}_{\min}\}$$

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Energy reduction

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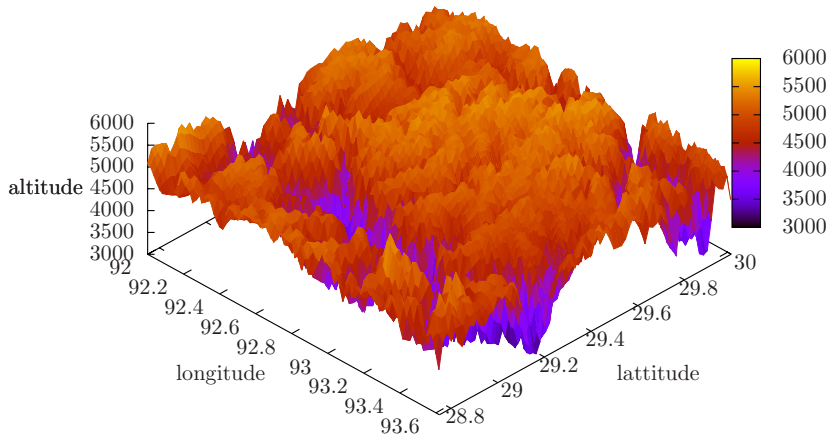
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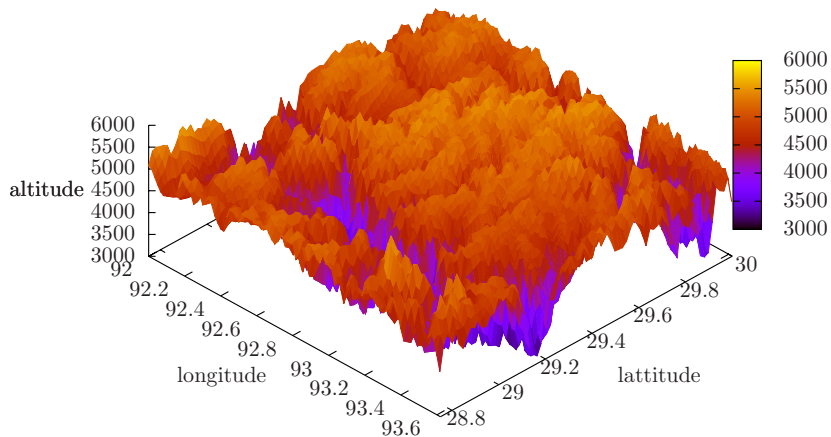
Landscape



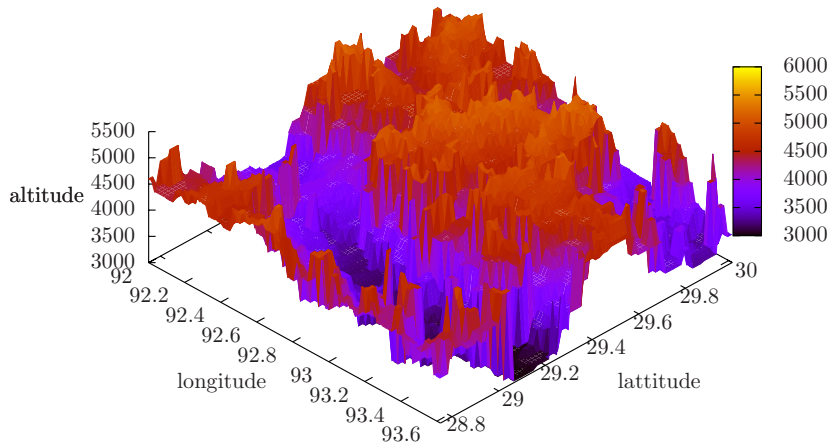
Reduced Landscape

$$E_{\text{red}}(\mathbf{S}) := E(\mathbf{S}_{\min}) \equiv E(G(\mathbf{S})).$$

Reduced Landscape



Reduced Landscape



Enable walk in reduced energy landscape

dynamical greedy algorithm

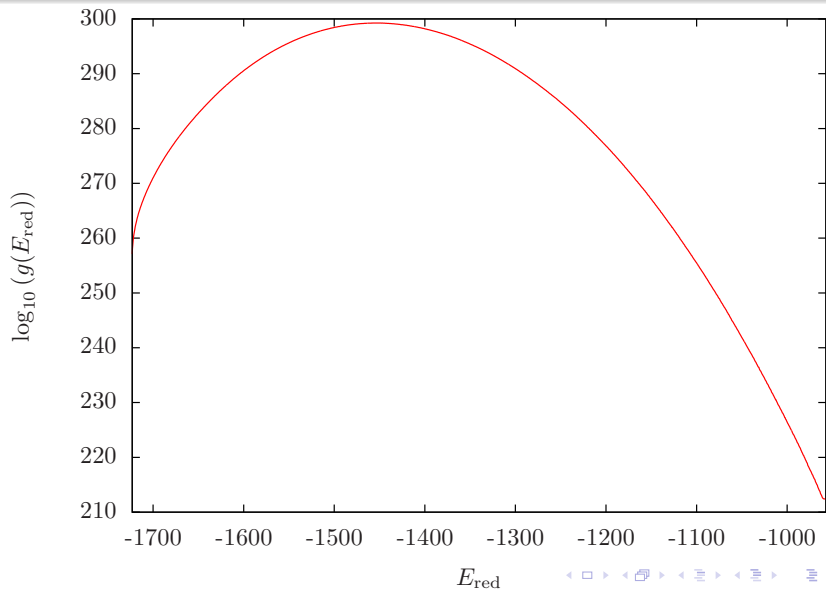
- Stores relevant information of a (standard) greedy optimization $\mathbf{S}_{\min} = G(\mathbf{S})$
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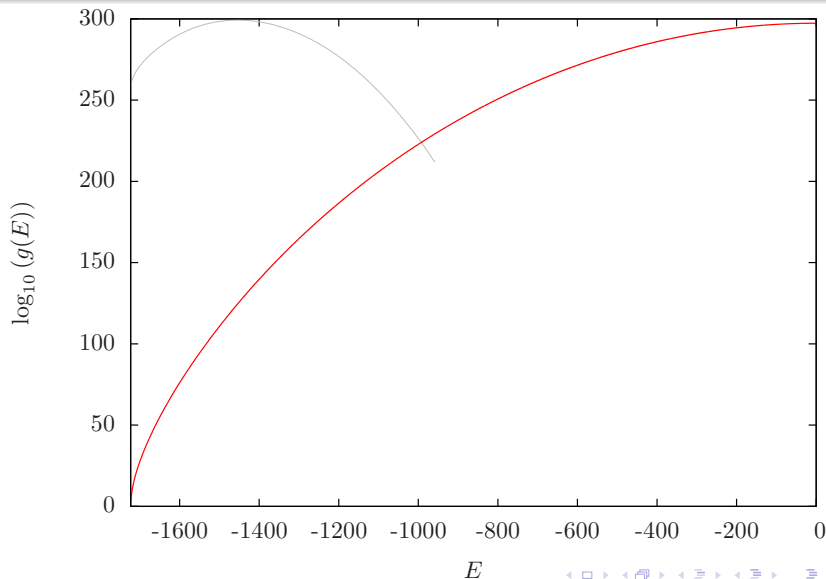
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Density of States as function of E_{red} (one sample, $L = 10$)



Density of States as function of E (one sample, $L = 10$)



Size of basin of attraction

- Perform a MC simulation within the basin.
(Reject all MC moves that change the result of the greedy optimization.)
- Relate the weight of the entire basin to the weight of the minimum configuration.
- E.g., flat-histogram over the overlap Q between current configuration \mathbf{S} and local minimum \mathbf{S}_{red} .

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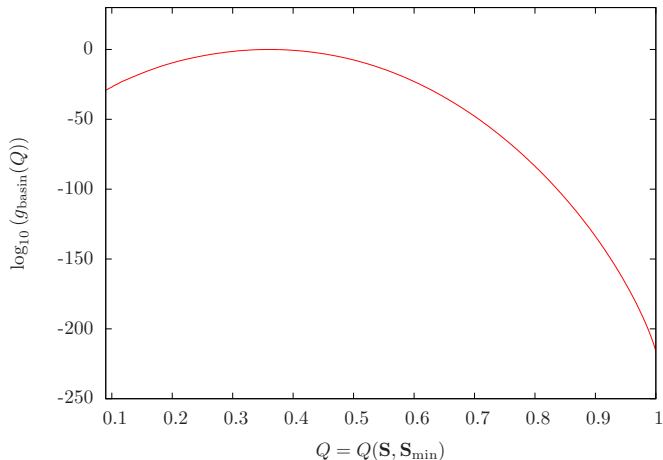
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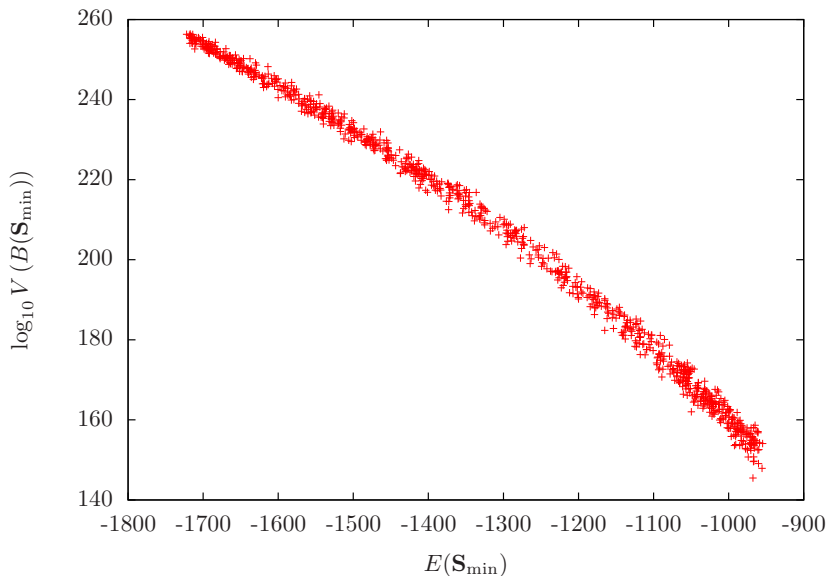
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Size of basin of attraction

- flat-histogram over the overlap Q between current configuration \mathbf{S} and local minimum \mathbf{S}_{\min}



Sizes of various basins



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- define logarithmic size : $s = \ln V$
- the sampled distribution $P^s(s)$ is biased towards large basins
- true distribution: $P_{\text{true}}^s(s) \propto \frac{1}{V} P^s(s) = e^{-s} P^s(s)$
- true distribution of sizes:
$$P_{\text{true}}^V(V) \propto \frac{1}{V^2} P^s(\ln V) = e^{-2s} P^s(s)$$
- $\langle V \rangle = \frac{\int_{-\infty}^{\infty} P^s(s) ds}{\int_{-\infty}^{\infty} e^{-s} P^s(s) ds}$
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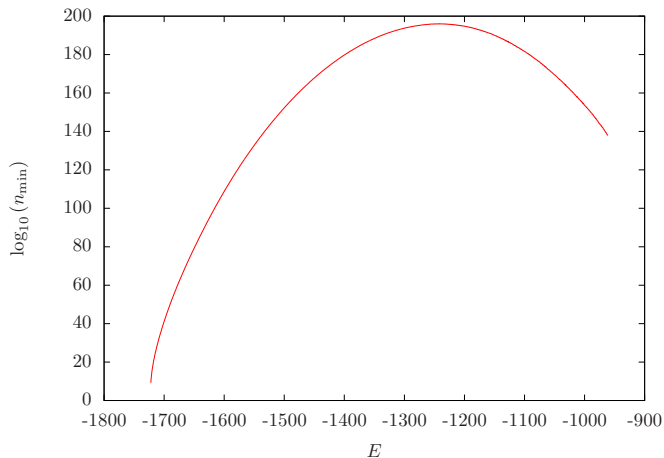
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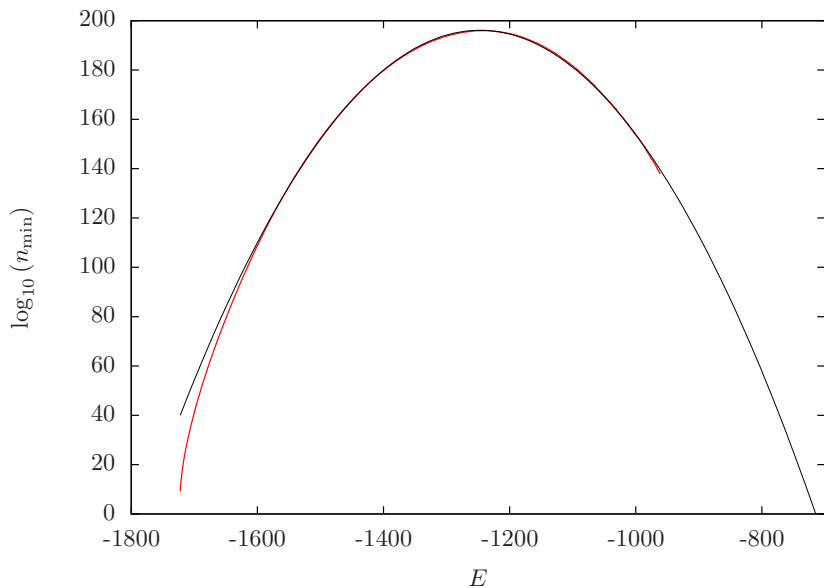
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Number of local minima $n_{\min} = g(E_{\text{red}})/\langle V \rangle(E_{\text{red}})$



Number of local minima



Concluding remarks

- Number and distribution of local minima can be determined.
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Thanks for your attention

Density of states fit

